

Derandomized Multi-block sign selection for PMEPR reduction of FBMC waveform

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Outline

The PAPR problem

FBMC signal model

Derandomized sign selection

Performance and simulations

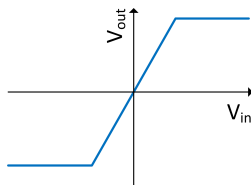
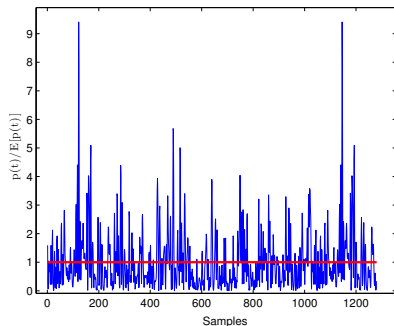
Conclusion

The PAPR problem

- ▶ High fluctuations of Multicarrier waveforms

$$s(n) = \sum_{k=0}^{N-1} X_k e^{j \frac{2\pi}{N} kn}$$

- ▶ Nonlinearity of Power Amplifier (PA)
- ▶ In-band and out-of-band distortion

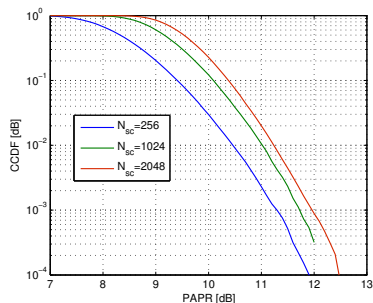


Definition and representation

- ▶ PMEPR for each symbol

$$\gamma = \frac{\max_{0 \leq t < T} |s(t)|^2}{\mathbb{E}[|s(t)|^2]}$$

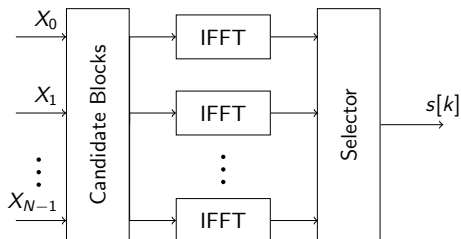
- ▶ Over intervals of length T , OFDM and OFDM/OQAM
- ▶ Complementary Cumulative Distribution Function (CCDF)



- ▶ $N=256$, 11.3 dB
- ▶ $N=1024$, 11.7 dB
- ▶ $N=2048$, 12 dB

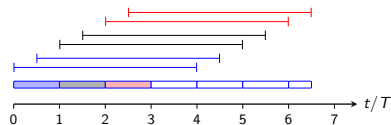
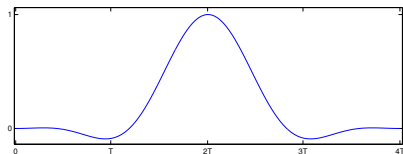
PAPR Reduction

- ▶ Many methods since 90s.
- ▶ Selected Mapping (SLM)
- ▶ $\mathbf{X} = [X_0, X_1, \dots, X_{N-1}]^T$
- ▶ $\mathbf{V}_m = [e^{j\phi_0}, e^{j\phi_1}, \dots, e^{j\phi_{N-1}}] \quad 1 \leq m \leq M$
- ▶ Sign change



FBMC signal model I

- ▶ A multicarrier modulation
- ▶ Flexible pulse shape
- ▶ $TF = \frac{1}{2}$ for real data symbols
- ▶ Applied 2x faster
- ▶ $TF = 1$ for complex data symbols
- ▶ component length $(K + \frac{1}{2})LN$



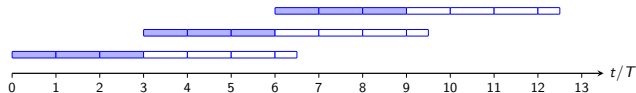
FBMC signal model II

- ▶ Continuous-time model:

$$s(t) = \sum_{m=-\infty}^{+\infty} \sum_{n=0}^{N-1} a_{n,m} h(t - m\tau_0) e^{j\frac{2\pi}{T}nt} e^{j\phi_{n,m}}$$

- ▶ a segment $P = QLN$

$$s[p] = s^{(r)}[p] + \sum_{m=0}^{2Q-1} \sum_{n=0}^{N-1} \epsilon_{q(n,m)} a_{n,m} \gamma_{n,m}(p) \quad 0 \leq p < LQN$$



FBMC signal model III

- ▶ Focus:

$$s[p] = s^{(r)}[p] + \sum_{m=0}^{2Q-1} \sum_{n=0}^{N-1} \epsilon_{q(n,m)} a_{n,m} \gamma_{n,m}(p) \quad 0 \leq p < LQN$$

- ▶ $\max |s[p]| \leq \max |s_r[p]| + \max |s_i[p]|$

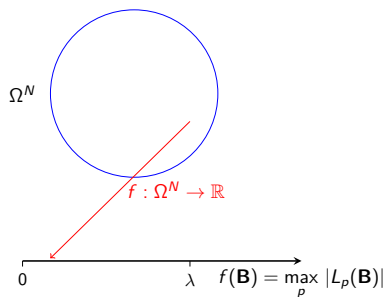
$$L_p(\epsilon) = L_p^r + \sum_{m=0}^{Q-1} \sum_{n=0}^{N-1} \epsilon_{q(n,m)} \Gamma_{n,m}(p) \quad 0 \leq p < 2P$$

- ▶ with $q(n, m) = (m - 1)N + n$

$$L_p = L_p^r + \sum_{\text{all}} \begin{bmatrix} \epsilon_{1,1} \Gamma_{0,0}(p) & \epsilon_{1,2} \Gamma_{0,1}(p) & \cdots & \epsilon_{1,Q} \Gamma_{0,Q-1}(p) \\ \epsilon_{2,1} \Gamma_{1,0}(p) & \epsilon_{2,2} \Gamma_{1,1}(p) & \cdots & \epsilon_{2,Q} \Gamma_{1,Q-1}(p) \\ \vdots & \vdots & \ddots & \vdots \\ \epsilon_{N,1} \Gamma_{N-1,0}(p) & \epsilon_{N,2} \Gamma_{N-1,1}(p) & \cdots & \epsilon_{N,Q} \Gamma_{N-1,Q-1}(p) \end{bmatrix}$$

Derandomized Algorithm I

- ▶ A random vector $\mathbf{B} \in \Omega^N$,
Equiprobable, $\Omega = \{-1, 1\}$,
 $\mathbf{B} = [B_1, B_2, \dots, B_N]^T$
- ▶ BAD event
 $A = \{\mathbf{B} : f(\mathbf{B}) \geq \lambda\}$
- ▶ $P(A) < 1 \rightarrow$ Exists B 's with A
not true



- ▶ Choose signs b_j iteratively such that

$$1 > P(A) \geq P(A|b_{1:j-1} = b_{1:j-1}^*) \geq P(A|b_{1:j} = b_{1:j}^*) \geq P(A|b_{1:N} = b_{1:N}^*) = 0$$

- ▶ Equiprobable RV \rightarrow

$$\begin{aligned} P(A|b_{1:j-1} = b_{1:j-1}^*) &= \sum_{\theta \in \Omega} P(\theta) P(A|b_{1:j-1} = b_{1:j-1}^*, b_j = \theta) \\ &\geq \min_{\theta \in \Omega} P(A|b_{1:j-1} = b_{1:j-1}^*, b_j = \theta) \end{aligned}$$

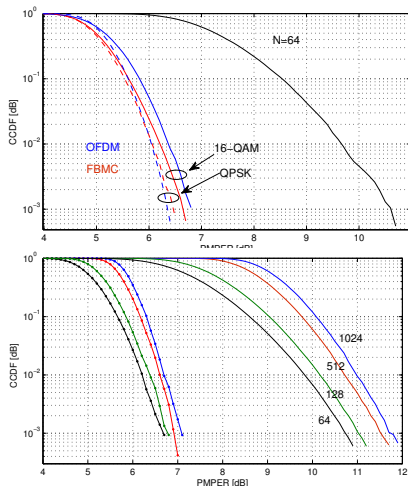
Derandomized Algorithm II

- ▶ Exhaustive search?!
- ▶ Upperbound $P(A|b_{1:j} = b_{1:j}^*) \leq F(b_{1:j}^*)$
 - ▶ (i) $F < 1$
 - ▶ (ii) $F(b_{1:j-1}^*) \geq \min_{\theta \in \Omega} F(b_{1:j-1}^*, b_j = \theta)$
- ▶ In this problem:
 - ▶ Union bound:
$$P\left(\max_{\theta \in \{-1,1\}} |L_p(\epsilon)| > \lambda | \epsilon_{1:j} = \epsilon_{1:j}^* \right) \leq \sum_p P(|L_p(\epsilon)| > \lambda | \epsilon_{1:j} = \epsilon_{1:j}^*)$$
 - ▶ Chernoff bound: $P(X > \mu) = P(e^{tX} > e^{t\mu}) < \frac{E[e^{tX}]}{e^{t\mu}}$
- ▶ Solution:
 - ▶ $F_p(\epsilon_{1:j}^*) = 2e^{-t\lambda} \cosh t(L_p(\epsilon_{1:j}^*) + L_p^r) \prod_{(n,m) \in \rho_j^c} \cosh t\Gamma_{n,m}(p)$
 - ▶ $\lambda = \sqrt{2(3QNE_{\max}^o h_{\max} + L_m^2) \ln 4P}$

Performance

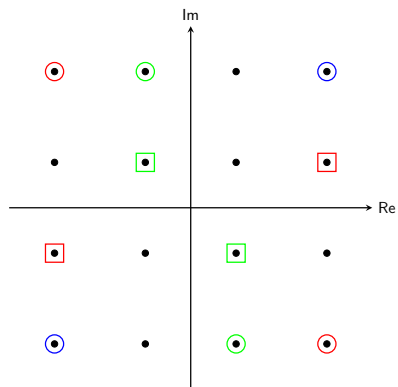
- ▶ Same for $Q = 1$ and higher.
- ▶ Slightly better for QPSK
- ▶ Almost the same for OFDM and FBMC
- ▶ PMPER grows faster for uncoded signal

N	at CCDF= 10^{-3}
64	4 dB
128	4.2 dB
512	4.6 dB
1024	4.6 dB



Receiver

- ▶ Side Information or Rate Loss
- ▶ Detection by discarding signs



Conclusion

- ▶ High reduction performance.
- ▶ Works as well with FBMC waveform and overlappings
- ▶ Same performance as that of OFDM.

Thank you!